

Braidings on Non-Split Tambara-Yamagami Categories over the Reals

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Topics

Introduction

- ▶ Categories
- ▶ Fusion Categories / Tambara-Yamagami Categories
- ▶ Split Case (Example Computation)
- ▶ Non-Split Case: Real-Quaternionic

Categories

Idea: Mathematical objects and maps between them

- ▶ Sets and functions
- ▶ Vector spaces and linear transformations
- ▶ Groups and group homomorphisms
- ▶ Rings and ring homomorphisms
- ▶ Topological spaces and continuous functions

Definition

Categories

objects	$a, b, c \in \mathcal{C}$
morphisms	$(a \xrightarrow{f} b) \in \text{Hom}(a, b)$
identity	$\exists (a \xrightarrow{\text{id}_a} a) \in \text{Hom}(a, a)$
composition	$\forall (a \xrightarrow{f} b) \forall (b \xrightarrow{g} c) \exists (a \xrightarrow{g \circ f} c)$
associativity	$\forall (a \xrightarrow{f} b \xrightarrow{g} c \xrightarrow{h} d) \quad (hg)f = h(gf)$
unitality	$\forall (a \xrightarrow{f} b) \implies f \text{id}_a = \text{id}_b f = f$

String Diagrams

Categories

<p>Objects: strings $\begin{array}{c} \\ a \end{array}$, $\begin{array}{c} \\ b \end{array}$, etc.</p>	<p>Morphisms: coupons $\begin{array}{c} b \\ \\ \boxed{f} \\ \\ a \end{array}$</p>
<p>Identity: empty string $\begin{array}{c} \\ a \end{array}$</p>	<p>Composition: vertical stacking $\begin{array}{c} c \\ \\ \boxed{g} \\ \\ \boxed{f} \\ \\ a \end{array}$</p>

String Diagrams (Continued)

Categories

Associativity: correctness of (no parenthesis)



Unitality: correctness of (can compose with empty string)



Fusion Categories

Idea: “categorified” rings

Ring R	Fusion Category \mathcal{C}
R set, elements $a, b, c \in R$	objects $A, B, C \in \text{ob}(\mathcal{C})$ morphisms $A \xrightarrow{f} C, B \xrightarrow{g} D$
sums $a + b$	direct sums $A \oplus B$
products $a \times b$	tensor products $A \otimes B, C \otimes D$ (functor) $A \otimes B \xrightarrow{f \otimes g} C \otimes D$
associativity $(ab)c = a(bc)$	<u>associators</u> $(A \otimes B) \otimes C \xrightarrow[\cong]{\alpha_{A,B,C}} A \otimes (B \otimes C)$

(Defn: A fusion category is a rigid finite semisimple linear monoidal category.)

Pentagon Axiom

Fusion Categories

Q. How to tensor four objects?

(In rings, $((ab)c)d = (a(bc))d = a((bc)d) = a(b(cd)) = (ab)(cd) =: abcd$)

In fusion categories, two ways to go between $((AB)C)D$ to $A(B(CD))$
(omitting \otimes s):

(continued on next slide)

Pentagon Axiom

Fusion Categories

Q. How to tensor four objects?

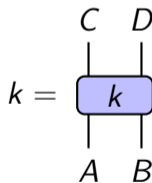
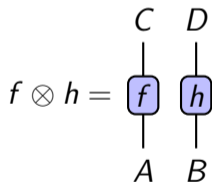
$$\begin{array}{ccc}
 ((AB)C)D & \xrightarrow{\alpha_{AB,C,D}} & (AB)(CD) \\
 \alpha_{A,B,C} \otimes \text{id}_D \downarrow & & \searrow \alpha_{A,B,CD} \\
 (A(BC))D & & A(B(CD)) \\
 \searrow \alpha_{A,BC,D} & & \nearrow \text{id}_A \otimes \alpha_{B,C,D} \\
 & A((BC)D) &
 \end{array}$$

We require them to be the same (diagram commutes).

String Diagrams (Con't)

Fusion Categories

We denote tensor products on string diagrams by stacking boxes horizontally:
 $\forall (A \xrightarrow{f} C), \forall (B \xrightarrow{h} D), \forall (A \otimes B \xrightarrow{k} C \otimes D)$, we write



Braidings

Fusion Categories

Idea: Categorifying commutativity

Definition: A *braiding* on a fusion category \mathcal{C} is a (natural) set of isomorphisms $c_{X,Y}: X \otimes Y \xrightarrow{\cong} Y \otimes X$ for all objects $X, Y \in \mathcal{C}$, denoted as

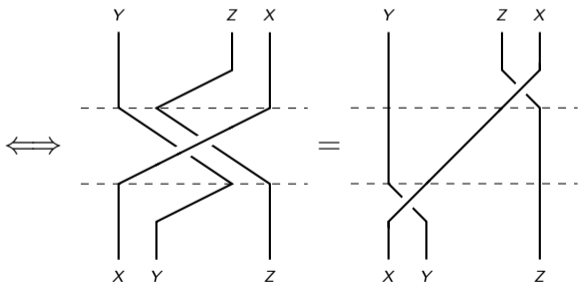
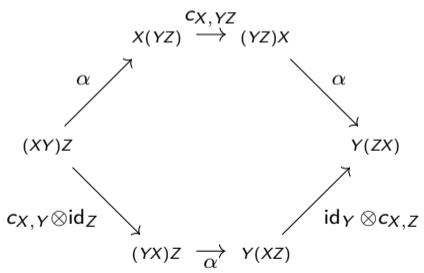


satisfying the following conditions: (continued on next slide)

Braidings

Fusion Categories

Definition: (continued) 1) hexagon axiom

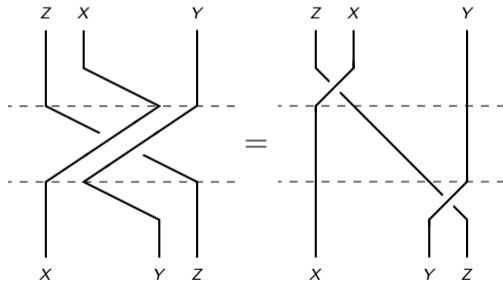


Braidings

Fusion Categories

and 2) inverse hexagon axiom

$$\begin{array}{ccccc}
 & & (XY)Z \xrightarrow{c_{XY,Z}} & Z(XY) & \\
 \alpha^{-1} \nearrow & & & & \searrow \alpha^{-1} \\
 X(YZ) & & & & (ZX)Y \\
 \downarrow \text{id}_X \otimes c_{Y,Z} & & & & \uparrow c_{X,Z} \otimes \text{id}_Y \\
 X(ZY) \xrightarrow{\alpha^{-1}} & (XZ)Y & & &
 \end{array}$$

 \Leftrightarrow


Tambara-Yamagami Categories

Fusion Categories

Idea: Categorifying “near-group” rings

Let G be a finite group. The (*split*) Tambara-Yamagami fusion category has $G \sqcup \{m\}$ as a set of simple objects (“basis”), and objects are generated under finite direct sums. The tensor product is defined on simples: (for all $a, b \in G$)

$$a \otimes b = ab \quad , \quad a \otimes m = m = m \otimes a \quad , \quad m \otimes m = \bigoplus_{c \in G} c,$$

as \otimes distributes over \oplus (these are called “fusion rules”).

Split vs. Non-Split

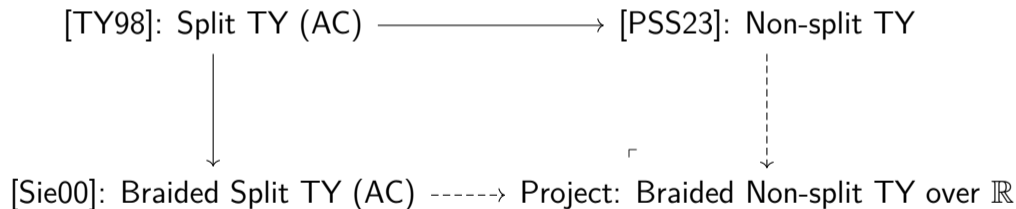
Fusion Categories

(Helpful fact: fusion categories are linear, i.e. Hom's are vector spaces over \mathbb{F} .)

Definition: A simple object $X \in \mathcal{C}$ is *split* if $\text{End}(X) := \text{Hom}(X, X) \cong \mathbb{F}$. Otherwise, it is non-split. A category is split if all of its simple objects are.

Project

Our project aims to construct this “pushout” diagram:



Strategy, briefly

1. Pick “basis” for category: objects are finite direct sums of simple objects; morphisms look like “matrices”
2. Schur’s lemma: tells me what morphisms are between simple objects
3. Yoneda lemma: examine morphisms by looking at effects of precomposition, which are *linear maps* between Hom spaces
4. Pick bases for Hom spaces and compute

Classification of Associators

Split Case

Theorem: [Tambara-Yamagami '98] Any split TY category is determined by a triple (G, χ, τ) , where G is a finite group, $\chi: G \times G \rightarrow \mathbb{R}^\times$ is a nondegenerate symmetric bicharacter, and $\tau \in \{\pm 1/\sqrt{|G|}\}$. Associators can be explicitly constructed:

$$\alpha_{m,a,m} = \begin{array}{c} | \quad | \\ | \quad | \\ | \quad | \\ | \quad | \\ m \quad a \quad m \end{array} = \sum_{d \in G} \chi(a, d) \cdot \left(\begin{array}{c} m \quad a \quad m \\ \diagdown \quad \diagup \\ \quad \quad d \\ \diagup \quad \diagdown \\ m \quad a \quad m \end{array} \right)$$

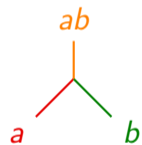
The diagram on the left shows three vertical lines labeled m , a , and m at the bottom. The middle line a has a red line that starts at the bottom, goes up, then right, then up again to meet the top of the a line.

The diagram on the right is a crossing of two lines labeled a at the bottom, with a vertical line labeled d in the middle. The top of the crossing is labeled with m , a , and m from left to right. The bottom of the crossing is also labeled with m , a , and m from left to right. The lines are black, except for the a lines which are red.

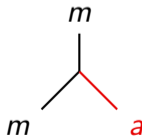
Bases for Hom's

Split Case

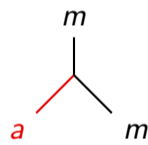
(continued) where trivalent vertices denote a set of fixed basis vectors:



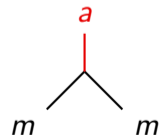
$$\in \text{Hom}(a \otimes b, ab)$$



$$\in \text{Hom}(m \otimes a, m)$$



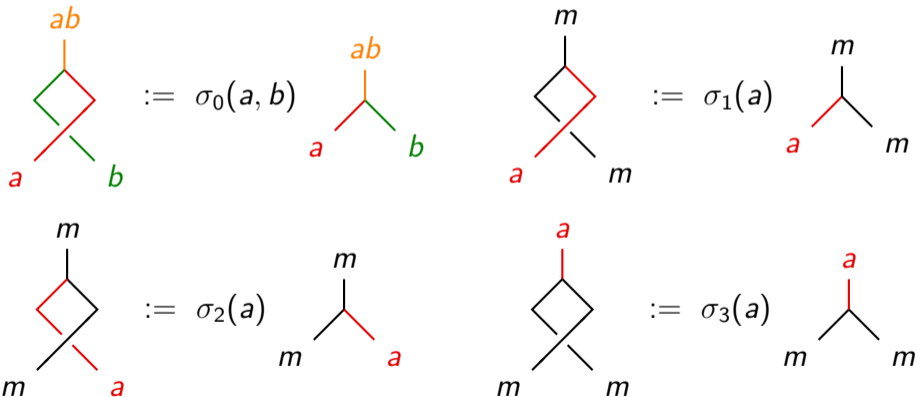
$$\in \text{Hom}(a \otimes m, m)$$



$$\in \text{Hom}(m \otimes m, a)$$

Braiding Coefficients

Split Case



Example Computation

Split Case

(Top path, hexagon 6)

$$\begin{array}{c} b \\ | \\ m \quad a \quad m \end{array} \xrightarrow{\alpha_{m,a,m}^*} \sum_{d \in A} \chi(a, d) \cdot \left(\begin{array}{c} b \\ \diagdown \quad \diagup \\ m \quad a \quad m \\ | \\ d \\ \diagdown \quad \diagup \\ m \quad a \quad m \end{array} \right)$$

Example Computation

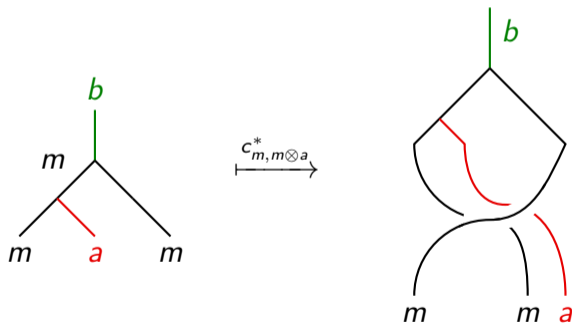
Split Case

(Top path, hexagon 6)

$$\begin{array}{c} b \\ | \\ m \quad \quad m \\ / \quad \backslash \\ m \quad a \quad m \end{array} \xrightarrow{\alpha_{m,a,m}^*} \chi(a, b) \cdot \left(\begin{array}{c} b \\ | \\ m \quad \quad m \\ / \quad \backslash \\ m \quad a \quad m \end{array} \right)$$

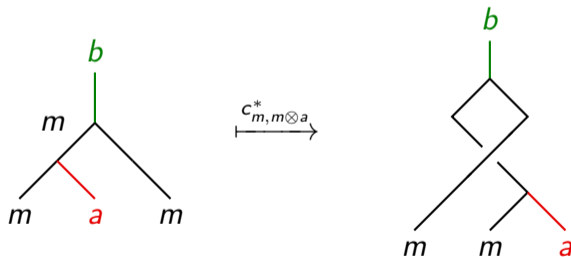
Example Computation

Split Case



Example Computation

Split Case



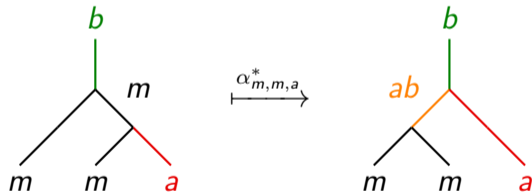
Example Computation

Split Case

$$\begin{array}{c} b \\ \diagdown \quad \diagup \\ \diagup \quad \diagdown \\ m \quad m \quad a \end{array} = \sigma_3(b) \cdot \left(\begin{array}{c} b \\ \diagdown \quad \diagup \\ m \quad m \\ \diagdown \quad \diagup \\ m \quad m \quad a \end{array} \right)$$

Example Computation

Split Case



Example Computation

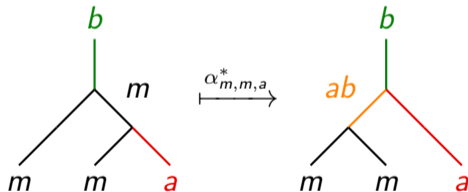
Split Case

(Bottom path, hexagon 6)

$$\begin{array}{c} b \\ | \\ \diagdown \quad \diagup \\ m \quad a \quad m \end{array} \xrightarrow{(\text{id}_m \otimes c_{a,m})^*} \left(\begin{array}{c} b \\ | \\ \diagdown \quad \diagup \\ m \quad m \quad a \end{array} \right) = \sigma_2(a) \cdot \left(\begin{array}{c} b \\ | \\ \diagdown \quad \diagup \\ m \quad m \quad a \end{array} \right)$$

Example Computation

Split Case



Example Computation

Split Case

$$\begin{array}{c}
 \begin{array}{c}
 b \\
 | \\
 ab \\
 / \quad \backslash \\
 m \quad m \\
 \backslash \quad / \\
 \quad m
 \end{array}
 \quad a
 \end{array}
 \xrightarrow{(c_{m,m} \otimes \text{id}_a)^*}
 \left(
 \begin{array}{c}
 b \\
 | \\
 ab \\
 / \quad \backslash \\
 m \quad m \\
 \backslash \quad / \\
 \quad m
 \end{array}
 \quad a
 \right)
 =
 \left(
 \sigma_3(ab) \cdot
 \begin{array}{c}
 b \\
 | \\
 ab \\
 / \quad \backslash \\
 m \quad m \\
 \backslash \quad / \\
 \quad m
 \end{array}
 \quad a
 \right)$$

Example Computation - Result

Split Case

$$\sigma_3(b)\chi(a, b) \left(\begin{array}{c} b \\ | \\ ab \\ / \quad \backslash \\ m \quad m \quad a \end{array} \right) = \sigma_3(ba)\sigma_2(a) \left(\begin{array}{c} b \\ | \\ ab \\ / \quad \backslash \\ m \quad m \quad a \end{array} \right)$$

Linear independence:

$$\sigma_3(b)\chi(a, b) = \sigma_3(ba)\sigma_2(a)$$

Do this fifteen more times to get: (con't on next slide)

Hexagon Equations

Split Case

Hexagon axioms \iff sixteen hexagon equations \iff six reduced equations:

$$\sigma_0(a, b) = \chi(a, b), \quad (1)$$

$$\sigma_1(a)^2 = \chi(a, a), \quad (2)$$

$$\sigma_1(ab) = \sigma_1(a)\sigma_1(b)\chi(a, b), \quad (3)$$

$$\sigma_2(a) = \sigma_1(a), \quad (4)$$

$$\sigma_3(1)^2 = \tau \sum_{c \in G} \sigma_1(c), \quad (5)$$

$$\sigma_3(a) = \sigma_3(1)\sigma_1(a)\chi(a, a). \quad (6)$$

Equation (2) tells us that $\chi(a, a) > 0$ for all $a \in G$, since $\sqrt{-1} \notin \mathbb{R}$.

Classification of Braiding

Split Case

Equation (2) tells us that $\chi(a, a) > 0$ for all $a \in G$, which places a big restriction on χ ([Wall '63]), leading to the following classification.

Theorem: [J.] Any split TY category over \mathbb{R} that admits a braiding is equivalent to $\mathcal{C}(K_4^n, h^{\oplus n}, \tau)$, where K_4 denotes the Klein four-group, h denotes the hyperbolic pairing on K_4 , $\tau \in \{\pm 1/2^n\}$, and $n \in \mathbb{N}$.

Classification of Braidings

Split Case

If we also consider the restriction

$$\sigma_3(1)^2 = \tau \sum_{c \in G} \sigma_1(c),$$

we have a classification of the solutions to the hexagon equations (the braidings), as well as a classification up to (monoidal) equivalence:

There are two non-equivalent braidings on $\mathcal{C}(K_4^n, h^{\oplus n}, \tau)$.

Non-Split Cases

For all simples $X \in \mathcal{C}$, $\text{End}(X)$ is not necessarily the ground field \mathbb{R} , but we have the following:

- ▶ Lemma (Schur): $\text{End}(X)$ is a division algebra over \mathbb{R}
- ▶ Theorem (Frobenius): Finite dimensional division algebra over \mathbb{R} is \mathbb{R} , \mathbb{C} , or \mathbb{H}
- ▶ Results from [Plavnik-Sanford-Sconce '23]:
 - $\text{End}(g) \cong \text{End}(1_G)$ for all $g \in G$;
 - $\text{End}(1_G)$ is commutative;
 - $\text{End}(1_G) \subseteq Z(\text{End}(m))$

Non-Split Cases

$\text{End}(1_G)$	$\text{End}(m)$	Viable?
\mathbb{R}	\mathbb{R}	Split
	\mathbb{H}	Case 1
	\mathbb{C}	Case 2
\mathbb{C}	\mathbb{R}	$Z(\mathbb{R}) = \mathbb{R}$
	\mathbb{H}	$Z(\mathbb{H}) = \mathbb{R}$
	\mathbb{C}	Case 3

Case 3 does not admit braidings. We will look at Case 1 for the rest of the talk.

Classification of Associators

Real/Quaternionic Case

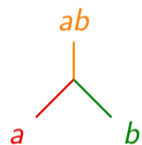
The fact that $Z(\mathbb{H}) = \mathbb{R}$ makes the real/quaternionic case feel very similar to the split case.

Theorem: [Plavnik-Sanford-Sconce '23] All real/quaternionic TY categories arise as $\mathcal{C}_{\mathbb{H}}(G, \chi, \tau)$, where G is a finite group, $\chi : G \times G \rightarrow \mathbb{R}^{\times}$ is a nondegenerate symmetric bicharacter, and $\tau \in \{\pm 1/\sqrt{4|G|}\}$.

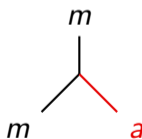
Associators can be constructed using fixed basis vectors: (continued on next slide)

Bases for Hom's

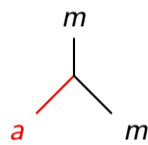
Real/Quaternionic Case



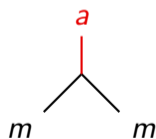
$$\in \text{Hom}(a \otimes b, ab)$$



$$\in \text{Hom}(m \otimes a, m)$$



$$\in \text{Hom}(a \otimes m, m)$$



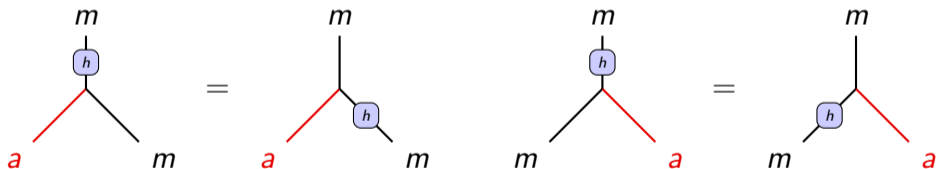
$$\in \text{Hom}(m \otimes m, a)$$

these have an additional nice property: (continued on next slide)

Bases for Hom's

Real/Quaternionic Case

The chosen basis vectors have the following property: we can pass all $h \in \mathbb{H}$ through vertices:

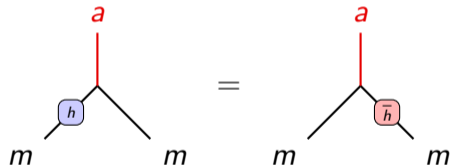


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Bases for Hom's

Real/Quaternionic Case

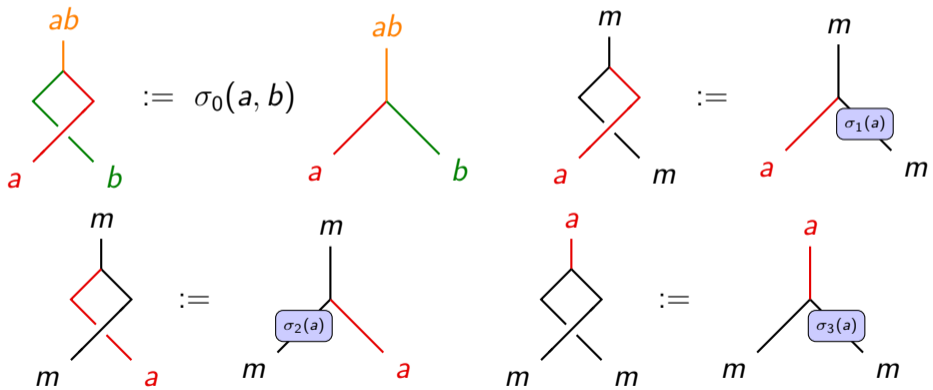
(continued)



... but sometimes they get conjugated.

Braiding Coefficients

Real/Quaternionic Case



Lemma

Non-Split Cases

...but the σ 's actually are all real valued!

Lemma: The function σ_1 is real-valued.

Proof: We want to show that for all $a \in A$ and all $h \in \mathbb{H}$,

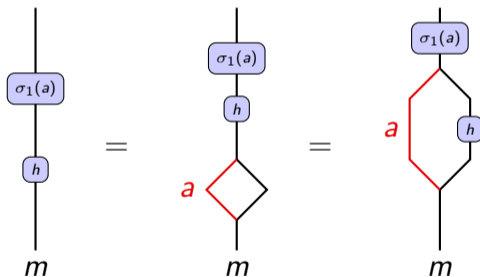
$$\sigma_1(a) \cdot h = h \cdot \sigma_1(a),$$

which would tell us that $\sigma_1(a) \in (Z(\mathbb{H}) = \mathbb{R})$. (con't on next slide)

Lemma (Proof)

Non-Split Cases

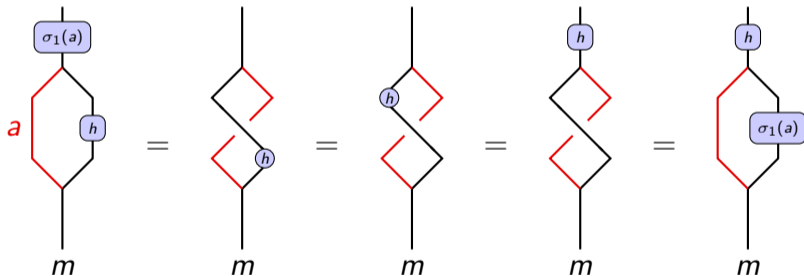
With graphical calculus, we have



Lemma (Proof)

Non-Split Cases

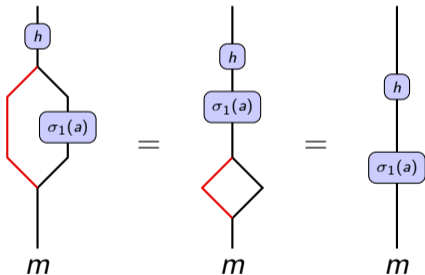
Now, by the definition of σ_1 and naturality, we have



Lemma (Proof)

Non-Split Cases

Pass $\sigma_1(a)$ through the trivalent vertex to get



as desired. □

Hexagon Equations

Real/Quaternionic Case

The reduced hexagon are as follows (note that (\star) is the only one different from the split case):

$$\sigma_0(a, b) = \chi(a, b),$$

$$\sigma_1(a)^2 = \chi(a, a),$$

$$\sigma_1(ab) = \sigma_1(a)\sigma_1(b)\chi(a, b),$$

$$\sigma_2(a) = \sigma_1(a),$$

$$\sigma_3(1)^2 = -2 \cdot \tau \sum_{c \in A} \sigma_1(c), \quad (\star)$$

$$\sigma_3(a) = \sigma_3(1)\sigma_1(a)\chi(a, a).$$

Classification of Braidings

Real/Quaternionic Case

The classification is very similar to the split case (differences arise in the actual construction of the braidings).

Theorem: [J.] Any real-quaternionic TY category that admits a braiding is equivalent to $\mathcal{C}_{\mathbb{H}}(K_4^n, h^{\oplus n}, \tau)$, where $\tau \in \{\pm 1/2^{n+1}\}$. There are two non-equivalent braidings on each of these categories.

Future Directions

- ▶ Case 2 (Real/Complex): graphical computations produce a set of different equations which involve complex conjugation
- ▶ Case 3 (Complex/Complex): does not admit braidings but can admit “ G -crossed braidings”

References

- [PSS23] Julia Plavnik, Sean Sanford, and Dalton Sconce. *Tambara-Yamagami Categories over the Reals: The Non-Split Case*. 2023. arXiv: 2303.17843 [math.QA].
- [Sie00] Jacob A. Siehler. *Braided Near-group Categories*. 2000. arXiv: math/0011037 [math.QA].
- [TY98] Daisuke Tambara and Shigeru Yamagami. “Tensor categories with fusion rules of self-duality for finite abelian groups”. In: *Journal of Algebra* 209.2 (1998), pp. 692–707. DOI: <https://doi.org/10.1006/jabr.1998.7558>.
- [Wal63] C. T. C. Wall. “Quadratic forms on finite groups, and related topics”. In: *Topology* 2 (1963), pp. 281–298. DOI: 10.1016/0040-9383(63)90012-0.

Thank you!